# Modern Encryption Standard version IV: (MES-IV) 

Asoke Nath ${ }^{1}$, Payel Pal ${ }^{2}$


#### Abstract

In the present paper the authors have introduced a new cryptographic method named as Modern Encryption System Standard version IV which is basically a symmetric key cryptographic method. Here the authors have used three different type of cryptographic methods .Those method are columnar transposition method, bit level generalized vernam cipher method with feedback and bit wise XOR operation. This system is the extension of MES-III and partly Bit level Encryption Standard (BLES) -II and III. BLESI, II, III and MES-I, II, III developed by Nath et al where MES-I, II, III mostly based on mainly byte level encryption method. BLES-I, II and III are based on mostly bit level encryption methods. In the present MES-IV method authors have tried to combine different bit level encryption to make the entire encryption system more secured. The introduction of feedback in bit level generalized vernam cipher method prevents attacks such as differential attack or plain text attack. The random key generator has been used to construct the keypad for vernam cipher method. The present encryption method is free from common attacks and it is almost impossible to break the present method without knowing the exact key and the methods. This encryption method will be used to encrypt password, short messages, financial data etc.


## Keywords

Plain text, Cipher text, Randomization, Columnar Transposition, Feedback

## 1. Introduction

Keeping secrets is not easy. In fact human tendency is such that when told that something is a secret and asked to keep it secret, people are actually quite eager to share that secret with everyone else.

[^0]In the early days of serious computing there was not a great deal of emphasis on security, because the systems in those days were proprietary or closed. The chances of someone getting an access to the information being exchanged were not very high.

As the minicomputers and microcomputers evolved in the 1970s and 1980s, the issue of information security started to gain more prominence. However it was the internet, which changed the whole computing paradigm and brought a tremendous change in the way computers communicated with each other. The world of computers had suddenly become very open. Therefore, it is very important to know how we can make information exchange secure. The confidentiality and genuineness of data has now become a very important issue. To send any important information from one user to another user normally the people are using e-mail as their transmission media. But the message of the e-mail can be trapped by the hacker between sender and receiver provided it is in raw form. To get rid of this problem one has to send the encrypted text or cipher text from client to server or to another client. To protect any kind of hacking problems nowadays network security and cryptography is an emerging research area where the programmers are trying to develop some strong encryption algorithm so that no intruder can intercept the encrypted key cryptography.

The cryptography methods can be divided into two categories: (i) symmetric key cryptography where one key is used for both encryption and decryption purpose. (ii) Public key cryptography where two different keys are used one for encryption and the other for decryption purpose. In symmetric key we have to maintain only one key and hence the key management is simple. In public key cryptography we maintain two keys one is public key which is known to everybody and that can be used for encryption purpose and there is another key called private key which is kept secret key and that is used for decryption purpose only.
Nath et al [1,2,3,4,5,6] developed different symmetric key cryptosystem. The advantage of symmetric key method is that key management is
very simple. Recently Nath et al developed cryptography method called Modern Encryption Standard version-I and Modern Encryption Standard version-II. and Modern Encryption Standard version-III. The present version is called Modern Encryption Standard version-IV. Here the authors have used three independent cryptography methods namely columnar transposition method, bit-wise generalized vernam cipher method and bit wise XOR method. In the present version the different method can be used for large size of file encryption. The output shows that the encryption is very strong as the encrypted text is totally different if there is only one character different in two patterns. The standard encryption algorithm like RSA or DES if we apply on a pattern where all characters are same then after encryption the encrypted text will also show same repeated pattern. But the present method applied on repeated pattern but the output contains totally different pattern. In the present work the authors are proposing a symmetric key cryptography where firstly the keygen() function is called to generate the encryption number and the randomization number. Then the bit level columnar transposition encryption is applied block wise. After that bit level generalized vernam cipher method with feedback is applied. Hare the encryption key used is taken from the randomized two dimensional array. Finally the bit wise xor encryption is applied. This system of encryption can suitable to encrypt different financial secret data, corporate data, password, defense network etc.

## 2. Algorithm

## Bitwise Columner Transposition Encryption Function Algorithm : col_trans_enc(file f1,file f2)

## This function takes two files f1 and f2 as argument

| step | 1 | $:$ | Open f1 in read mode. |
| :--- | :--- | :--- | :--- |
| step | 2 | $:$ | Open a file f 2 in write mode. |
| step | 3 | $:$ | set $\mathrm{i}=0$ and $\mathrm{n}=0$ |
| step | 4 | $:$ | set $\mathrm{j}=0$ |
| step | 5 | $:$ | set $\operatorname{mat}[\mathrm{i}][\mathrm{j}]=\mathrm{n}$ |
| step | 6 | $:$ | set $\mathrm{n}=\mathrm{n}+1$ |
| step | 7 | $:$ | Set $\mathrm{j}=\mathrm{j}+1$ and $\mathrm{if} \mathrm{j}<16$ go to step 5 |
| step | 8 | $:$ | set $\mathrm{i}=\mathrm{i}+1$ and if $\mathrm{i}<16$ go to step 4 |
| step | 9 | $:$ | Call randomization () |
| step | 10 | $:$ | set $\mathrm{i}=0$ and set $\mathrm{j}=0$ |


| step <br> step | $\begin{aligned} & 11 \\ & 12 \end{aligned}$ | Go to the start of $f 1$ read next 256 character from f 1 and assign number of available character to n and assign the read characters to a[256] |
| :---: | :---: | :---: |
| step | 13 | if $\mathrm{n}<=0$ go to step 56 |
| step | 14 | set l=0 |
| step | 15 | if $1>=$ n go to step 22 |
| step | 16 | Call function char_to_bit(a[1], binary) //this assign each bit of ch to elements of array binary[8] |
| step | 17 | Set k=0 |
| step | 18 | if $\mathrm{k}>=8$ go to step 21 |
| step | 19 | Set table[1][k]=binary[k] +48 |
| step | 20 | Set $\mathrm{k}=\mathrm{k}+1$ and go to step 18 |
| step | 21 | Set $\mathrm{l}=1+1$ and go to step 15 |
| step | 22 | assign all element of filearr[8] to 0 |
| step | 23 | set $\mathrm{il}=0$ and $\mathrm{j} 1=0$ |
| step | 24 | set $\mathrm{k}=0$ |
| step | 25 | if $\mathrm{k}>=8$ go to step 44 |
| step | 26 | Set temp=mat[i][j]\%8+1//a\%b returns the remainder after dividing a by b |
| step | 27 | set $\mathrm{i}=\mathrm{i}+1$ |
| step | 28 | If $\mathrm{i}=16$ set $\mathrm{i}=0$ and $\mathrm{j}=\mathrm{j}+1$ |
| step | 29 | If $\mathrm{j}=16$ set $\mathrm{j}=0$. |
| step | 30 | set m=0 |
| step | 31 | if $\mathrm{m}>=\mathrm{k}$ go to step 34 |
| step | 32 | if temp=filearr[m] go to step 34 |
| step | 33 | increase m by 1 and go to step 31 |
| step | 34 | if $\mathrm{m}<\mathrm{k}$ go to step 26 |
| step | 35 | set filearr[k]=temp |
| step | 36 | Set k=k +1 |
| step | 37 | set l=0 |
| step | 38 | if 1>=n go to step 43 |
| step | 39 | set table1[i1][j1]=table[1][temp-1] |
| step | 40 | set $\mathrm{j} 1=\mathrm{j} 1+1$ |
| step | 41 | if $\mathrm{j} 1=8$ set $\mathrm{j} 1=0$ and $\mathrm{i} 1=\mathrm{i} 1+1$ |
| step | 42 | set l=1+1 and go to step 38 |
| step | 43 | Go to step 25 |
| step | 44 | set l=0 |
| step | 45 | if 1>=n go t step 53 |
| step | 46 | set add=0 |
| step | 47 | set $\mathrm{k}=0$ |
| step | 48 | if $\mathrm{k}>=8$ go to step 51 |
| step | 49 | ```set add=add+ (table1[l][k]- 48)*power(7-j)``` |
| step | 50 | set $\mathrm{k}=\mathrm{k}+1$ and go to step 48 |
| step | 51 | Write add to file f2 |
| step | 52 | set l=1+1 and go to step 45 |
| step | 53 | if $\mathrm{n}<256$ go to step 56 |
| step | 54 | read next 256 character from f1 and assign number of available character |


|  |  | to n and assign the read characters to a[256] |
| :---: | :---: | :---: |
| step | 55 | Go to step 13 |
| step | 56 | Close all the files. |
| step | 57 | stop |
|  |  | bitwise vernam encryption with feedback:vernambitenc(file input,file output) |
| step | 1 | set k=0 |
| step | 2 | set $\mathrm{i}=0$ |
| step | 3 | if $\mathrm{i}>=16$ go to step 10 |
| step | 4 | set $\mathrm{j}=0$ |
| step | 5 | if $\mathrm{j}>=16$ go to step 9 |
| step | 6 | set mat[i][j]=k |
| step | 7 | increase k by 1 |
| step | 8 | increase j by 1 and go to step 5 |
| step | 9 | set $\mathrm{i}=\mathrm{i}+1$ and go to step 3 |
| step | 10 | call randomization() |
| step | 11 | set $\mathrm{i}=\mathrm{j}=0$ |
| step | 12 | set cr1=0 |
| step | 13 | open input file as fpn |
| step | 14 | read next character from fpn and assign to ch |
| step | 15 | if eof is found go to step 36 |
| step | 16 | call char_to_bit(ch,bitpattern[8]) |
| step | 17 | call char_to_bit(mat[i][j],key_bit[8]) |
| step | 18 | set $\mathrm{i}=\mathrm{i}+1$ |
| step | 19 | if $\mathrm{i}=16$ set $\mathrm{i}=0$ and set $\mathrm{j}=\mathrm{j}+1$ |
| step | 20 | if $\mathrm{j}=16$ set $\mathrm{j}=0$ |
| step | 21 | set cr=(bitpattern[0]+key_bit[0]+cr1)\%2 |
| step | 22 | set cb[0]=cr1=cr |
| step | 23 | set $\mathrm{k}=1$ |
| step | 24 | if $\mathrm{k}>=8$ go to step 29 |
| step | 25 | set cr=(bitpattern[k]+key_bit[k]+cr1)\%2 |
| step | 26 | set $\mathrm{cb}[\mathrm{k}]=\mathrm{cr}$ |
| step | 27 | set crl=cr |
| step | 28 | set $\mathrm{k}=\mathrm{k}+1$ and go to step 24 |
| step | 29 | set add=0 |
| step | 30 | set k=0 |
| step | 31 | if $\mathrm{k}>=8$ go to step 34 |
| step | 32 | set add=add+cb[k]*power(7-k) |
| step | 33 | increase k by 1 and go to step 31 |
| step | 34 | write add to file f1 |
| step | 35 | go to step 14 |
| step | 36 | set k=0 |
| step | 37 | set i=0 |
| step | 38 | if $\mathrm{i}>=16$ go to step 45 |
| step | 39 | set $\mathrm{j}=0$ |
| step | 40 | if $\mathrm{j}>=16$ go to step 44 |
| step | 41 | set mat $[\mathrm{i}][\mathrm{j}]=\mathrm{k}$ |


| step | 42 | increase k by 1 |
| :---: | :---: | :---: |
| step | 43 | increase j by 1 and go to step 40 |
| step | 44 | set $\mathrm{i}=\mathrm{i}+1$ and go to step 38 |
| step | 45 | call randomization() |
| step | 46 | set $\mathrm{i}=\mathrm{j}=0$ |
| step | 47 | set crl=0 |
| step | 48 | read next character from f1 and assign to ch |
| step | 49 | if eof is found go to step 70 |
| step | 50 | call char_to_bit(ch,bitpattern[8]) |
| step | 51 | call char_to_bit(mat[i][j],key_bit[8]) |
| step | 52 | set $\mathrm{i}=\mathrm{i}+1$ |
| step | 53 | if $\mathrm{i}=16$ set $\mathrm{i}=0$ and set $\mathrm{j}=\mathrm{j}+1$ |
| step | 54 | if $\mathrm{j}=16$ set $\mathrm{j}=0$ |
| step | 55 | : set |
| step | 56 | $\begin{aligned} & \text { cr=(bitpattern[0]+key_bit[0]+cr1)\%2 } \\ & \text { set cb[0]=cr1=cr } \end{aligned}$ |
| step | 57 | : set k=1 |
| step | 58 | : if $\mathrm{k}>=8$ go to step 63 |
| step | 59 | : set |
| step | 60 | $\begin{aligned} & \text { cr=(bitpattern[k]+key_bit[k]+crl)\%2 } \\ & \text { set } \operatorname{cb}[\mathrm{k}]=\mathrm{cr} \end{aligned}$ |
| step | 61 | : set crl=cr |
| step | 62 | : set $\mathrm{k}=\mathrm{k}+1$ and go to step 58 |
| step | 63 | : set add=0 |
| step | 64 | : set k=0 |
| step | 65 | : if $\mathrm{k}>=8$ go to step 68 |
| step | 66 | : set add=add+cb[k]*power(7-k) |
| step | 67 | : increase k by 1 and go to step 65 |
| step | 68 | : write add to file f2 |
| step | 69 | : go to step 48 |
| step | 70 | : set k=0 |
| step | 71 | : set $\mathrm{i}=0$ |
| step | 72 | : if $\mathrm{i}>=16$ go to step 79 |
| step | 73 | : set $\mathrm{j}=0$ |
| step | 74 | : if $\mathrm{j}>=16$ go to step 78 |
| step | 75 | : set mat[i][j]=k |
| step | 76 | : increase k by 1 |
| step | 77 | increase j by 1 and go to step 5 |
| step | 78 | : set $\mathrm{i}=\mathrm{i}+1$ and go to step 72 |
| step | 79 | : call randomization() |
| step | 80 | : set $\mathrm{i}=\mathrm{j}=0$ |
| step | 81 | : set cr1=0 |
| step | 82 | : read next character from f2 and assign to ch |
| step | 83 | : if eof is found go to step 104 |
| step | 84 | call char_to_bit(ch,bitpattern[8]) |
| step | 85 | call char_to_bit(mat[i][j],key_bit[8]) |
| step | 86 | set $\mathrm{i}=\mathrm{i}+1$ |
| step | 87 | : if $\mathrm{i}=16$ set $\mathrm{i}=0$ and set $\mathrm{j}=\mathrm{j}+1$ |
| step | 88 | if $\mathrm{j}=16$ set $\mathrm{j}=0$ |
| step | 89 | : set |
|  |  | cr=(bitpattern[0]+key_bit[0]+cr1)\%2 |


| step | 90 | set $\mathrm{cb}[0]=\mathrm{cr} 1=\mathrm{cr}$ |
| :---: | :---: | :---: |
| step | 91 | set k=1 |
| step | 92 | if $\mathrm{k}>=8$ go to step 97 |
| step | 93 | set |
| step | 94 | cr=(bitpattern[k]+key_bit[k]+cr1)\%2 set cb[k]=cr |
| step | 95 | set cr1=cr |
| step | 96 | set $\mathrm{k}=\mathrm{k}+1$ and go to step 92 |
| step | 97 | set add=0 |
| step | 98 | set k=0 |
| step | 99 | if $\mathrm{k}>=8$ go to step 102 |
| step | 100 | set add=add+cb[k]*power(7-k) |
| step | 101 | increase k by 1 and go to step 99 |
| step | 102 | write add to file f3 |
| step | 103 | go to step 82 |
| step | 104 | call file_rev(f3,f4) |
| step | 105 | set $\mathrm{k}=0$ |
| step | 106 | set $\mathrm{i}=0$ |
| step | 107 | if $\mathrm{i}>=16$ go to step 114 |
| step | 108 | set $\mathrm{j}=0$ |
| step | 109 | if $\mathrm{j}>=16$ go to step 113 |
| step | 110 | set mat[i][j]=k |
| step | 111 | increase k by 1 |
| step | 112 | increase j by 1 and go to step 109 |
| step | 113 | set $\mathrm{i}=\mathrm{i}+1$ and go to step 107 |
| step | 114 | call randomization() |
| step | 115 | set $\mathrm{i}=\mathrm{j}=0$ |
| step | 116 | set cr1=0 |
| step | 117 | read next character from f 4 and assign to ch |
| step | 118 | if eof is found go to step 139 |
| step | 119 | call char_to_bit(ch,bitpattern[8]) |
| step | 120 | call char_to_bit(mat[i][j],key_bit[8]) |
| step | 121 | set $\mathrm{i}=\mathrm{i}+1$ |
| step | 122 | if $\mathrm{i}=16$ set $\mathrm{i}=0$ and set $\mathrm{j}=\mathrm{j}+1$ |
| step | 123 | if $\mathrm{j}=16$ set $\mathrm{j}=0$ |
| step | 124 | set cr=(bitpattern[0]+key_bit[0]+cr1)\%2 |
| step | 125 | set cb[0]=crl=cr |
| step | 126 | set $\mathrm{k}=1$ |
| step | 127 | if $\mathrm{k}>=8$ go to step 132 |
| step | 128 | set cr=(bitpattern[k]+key_bit[k]+cr1)\%2 |
| step | 129 | set cb[k]=cr |
| step | 130 | set cr1=cr |
| step | 131 | set $\mathrm{k}=\mathrm{k}+1$ and go to step 127 |
| step | 132 | set add=0 |
| step | 133 | set $\mathrm{k}=0$ |
| step | 134 | if $\mathrm{k}>=8$ go to step 137 |
| step | 135 | set add=add+cb[k]*power(7-k) |
| step | 136 | : increase k by 1 and go to step 134 |
| step | 137 | : write add to file output file |
| step | 138 | : go to step 117 |


| step | 139 | $:$ | close all files |
| :--- | :--- | :--- | :--- |
| step | 140 | $:$ | stop |
|  |  |  |  |
|  |  |  |  |
|  |  |  |  |
|  |  | Fit Wise XOR Encryption |  |
|  |  |  |  |

## This function takes two files f1 and f2 as argument

| step | 1 | Open the file f 1 in read mode |
| :---: | :---: | :---: |
| step | 2 | Open the file f2 in write mode |
| step | 3 | set l=size of file f1 |
| step | 4 | set $\mathrm{n} 1=1 / 32$ |
| step | 5 | set $\mathrm{n} 1=1 \% 32 / / \mathrm{a} \% \mathrm{~b}$ returns the remainder after dividing a by b |
| step | 6 | : Go to the start of file f1 |
| step | 7 | : set $\mathrm{i}=0$ and $\mathrm{n}=0$ |
| step | 8 | set $\mathrm{j}=0$ |
| step | 9 | set mat[i][j] $=\mathrm{n}$ |
| step | 10 | set $\mathrm{n}=\mathrm{n}+1$ |
| step | 11 | Set $\mathrm{j}=\mathrm{j}+1$ and if $\mathrm{j}<16$ go to step 9 |
| step | 12 | set $\mathrm{i}=\mathrm{i}+1$ and if $\mathrm{i}<16$ go to step 8 |
| step | 13 | set $\mathrm{i}=1$ |
| step | 14 | if i>secure then go to step 17 |
| step | 15 | call randomization() |
| step | 16 | set $\mathrm{i}=\mathrm{i}+1$ and go to step 14 |
| step | 17 | set $\mathrm{i}=1$ |
| step | 18 | : if i>n1 then go to step 23 |
| step | 19 | Read next 32 character from file f1 and assign it to array data1[32] |
| step | 20 | Call bit_stream(data1[32]) |
| step | 21 | Call encrypt_bit() |
| step | 22 | set $\mathrm{i}=\mathrm{i}+1$ go to step 18 |
| step | 23 | if $\mathrm{n} 2=0$ go to step 30 |
| step | 24 | : set $\mathrm{i}=0$ |
| step | 25 | if $\mathrm{i}>=\mathrm{n} 2$ go to step 30 |
| step | 26 | Read next character fron file f1 and assign to data2[i] of array data2[32] |
| step | 27 | set <br> data2[i]=rshift_residual(data2[i],5) |
| step | 28 | : write data2[i] to file f2 |
| step | 29 | : set i=i+1 go to step 25 |
| step | 30 | close all files |

[^1]
## bitwise vernam decryption with feedback：vernambitdec（file input，file output）

This algorithm is the reverse of vernambitenc algorithm．

Bitwise Xor Decryption Function：
Bitxordec（File F1，File F2）
This algorithm is the reverse of bitxorenc algorithm．

## 3．Results and Discussion

The Modern Encryption Standard version IV （MES－IV）is applied on different type of text．This is applied on various repeated patterns such as 1024 numbers of ASCII 0 to ASCII 16．The corresponding frequency graphs are shown below． The graph shows the diverse nature of encrypted text even the input text is of repeated same character．



Fig－1：Frequency Graph of ASCII code ＇0＇，＇1＇，＇2，＇3＇

This encryption method applied to different type of text／patterns and shown below are the pairs of such different type of patterns and the corresponding cipher text．

Table－1：Some known plain texts vs．encrypted text

|  | ORIGINAL TEXT | ENCRYPTED TEXT |
| :---: | :---: | :---: |
| 1 | AAAAAAAAAAAAAA AAAAAAAAAAAAAA AAAAAAAAAAAAAA AAAAAAAAAAAAAA AAAAAAAA 64－As） | 程갔血 칠下画晞黽黯晰浤 됟奮．椆成口酸에脆数 稲 펄ㅇ．ど竟頞懇䛀 |
| 2 | BAAAAAAAAAAAAA AAAAAAAAAAAAAA AAAAAAAAAAAAAA AAAAAAAAAAAAAA AAAAAAAAA （B＋64 As） |  ¿ žfö r <br>  ápAÔì Jf（国u đà ${ }^{\text {TM／}}$ |


| 3 | AAAAAAAAAAAAAA AAAAAAAAAAAAAA AAAAAAAAAAAAAA AAAAAAAAAAAAAA AAAAAAAAB （ $64 \mathrm{As}+\mathrm{B}$ ） | CC＞－＞o mTŸð $\square e ̀ A d] « b$ Ôī ${ }^{-}$II ô r <br> ¡®VôYÈÿ6®Ô ${ }^{\circ}$ Ö̈̈M£Ûû̂â ápAÔì Jfa ðà ${ }^{\text {TM／}}$ |
| :---: | :---: | :---: |
| 4 | 00000000 | J i！＇ý |
| 5 | 11111111 | ＂J i！＇ý |
| 6 | 01010101 | －J i！＇ý |
| 7 | 1111111100000000 | ［0A¢RZ£ WÜ¢R $\ddagger \mathrm{p}$ |
| 8 | HE IS GOOD | ¥ÒWÎêo ${ }^{\circ} £ 5$ w |
| 9 | abcabcabcabc | ，$\square$ Då $x^{0} 5$ |
| 10 | HE IS GOON | ¥Đuiê ${ }^{\circ} \mathrm{E}$ ¢5w |
| 11 | CE IS GOON | ${ }^{\text {Àqu }}{ }^{\circledR}{ }^{\circ} \mathrm{E} £ 5 \mathrm{w}$ |

In the table below the plain text and the corresponding encrypted text are shown．The text of Sl ．No． $12 \& 13$ are exactly same except the fourth character．However the encrypted texts are quite different under the same encryption key．The Sl．No． $14 \& 15$ shows the same thing．

|  | $\begin{aligned} & \text { ORIGINAL } \\ & \text { TEXT } \end{aligned}$ | ENCRYPTED TEXT |
| :---: | :---: | :---: |
| 12 | The root of all modern programming languages is ALGOL，introduc ed in the early 1960s．ALGOL was the first computer language to use a block structure． Although it never became popular in USA，it was widely used in Europe．ALGOL gave the concept of structured programming to the computer science community． | ```Ç8 iØãÕ®y=g' - Ç\|ef@Eãl!2BÑWÈ \(\not \subset O ̈\) §@1/4è```  ```Û®'VNTÁ6Cçæ» Á [Â.RŽ3Ï' \(\neg\) E' \(\mu \mathrm{p}\) ðfK- ,pE\|EŽbÒY゙‘ž冖GÉ¥Fã s7‘aN€@ iÉq_kú \({ }^{\text {º }}\). H~Đ4```   ```\(\%\) \%èF \(\left.4 c z^{2} \mathrm{e}\right] \mathrm{mfE}\)...ó êf \(\ddagger \mu \mathrm{U} \mu\) ' ' \(): ~\)```  ```\(1 \mathrm{WW} f \mathrm{y}\) 目- \#包(ŠÉU̇•Ü(O)×《Ěšm>s \(\neg \mathrm{RU} ```D9Ü``` |


| 13 | The boot of all modern programming languages is ALGOL，introduc ed in the early 1960s．ALGOL was the first computer language to use a block structure． Although it never became popular in USA，it was widely used in Europe．ALGOL gave the concept of structured programming to the computer science community． | ```iæ(>kÄ ÔÁy<æÓ ø包\%å? ¡7BÅwR, tÇ® \(\neg \sim\) 'ïh~ P òË包MF\# Ú g\&ÎQÁ"c\} Æ rà ÍÀ \(23 / 4\) ?®Ï \(\mu\)-Í à, \%o \(03 / 4 \mathrm{I}, 1 / 405 \& @\) Žvò```  ```SE . häð-\{把wå W•Ø! \(\cdot\) Ö 2 œ尹 e...F Û/\{\& TM园yф ÉèG\%uQ/V3\%...ogÕ \(\AA\) ¢ \(f>\ddot{\mathrm{i}}\) f"```   ```!OM£y? \#鸟(ŠÉÙ•Ü(C)×《Ěšm>s \(\neg \mathrm{RU} ```D9Ü``` |
| :---: | :---: | :---: |
| 14 | Modern batch processing is generally tied into a timesharing system．In this system the programme and the data are typed into the computer via a timesharing terminal or a personal computer． | öÖO <br>  <br> $=\$ « 6 \mathrm{Y} 0<\mathrm{Z}+\mathrm{B}=@ \AA$＠ A 7 <br> $1 / 2 \square H-g u ́ Z a ̊ \pm-£ i$ <br> ＇J）X uHgÂ a＂； $\mathrm{Ob}^{-}$KÕ $>$İI <br> ＜＿æfø İZ»ú！ $9^{*} æ E ̀ p H$ Ea <br> ié1／2aâîs＇Ïœ＾úò é Ê＜ùðbz－，F <br> ，ôð国Ò $f$ ！øåòM§た＂ $\mathrm{N} \neg$ çh <br> ÿ－$\left\{\hat{1}\right.$ ùk ${ }^{2} \AA ̊ \cdot 0 f \tilde{\mathrm{~N}}$ <br> DÃ•©～ó $€^{\sim} 3$ » <br> I3iÜZÖœ回äUK \％${ }^{\circ}$ |
| 15 | Moderp patch processing is generally tied into a timesharing system．In this system the programme and the data are typed into the computer via a timesharing terminal or a personal computer． | ＇c $\square$ ÃkÇûëK－ ú01＇＠8á£ ${ }^{\text {тм }} \mid 2$ İH＇$^{\prime} \S ®$ §ul€ \｛b－ ÖC $\div .6 \_i \sim E ́ \times U^{-} H u ̂ \nexists j$ <br> $\neg 1 E ̇ S U ́ Z ̌ H o ̈ D o ̈ G a c @, 7$ SŽ® Õ？LŽ＞i gÃ ${ }^{\circledR 1}$ ÍÊǔŠ ${ }_{i}<$ ò̀èêh£ ＠øéCKîs＇Kß＾\｛＇Ü ZTLyõbnIIIff，ôð回Ò $f$ ！ øåòM§Æ＂Ñ $\urcorner$ çh ÿ－\｛î ùk <br>  I3iÜZÖœ国UK \％${ }^{\circ}$ |

## 4．Conclusion

The results above show how strong the encryption method is．Three different methods of encryption
techniques have been applied. The feedback mechanism gives to the strength of the encryption. The random matrix is generated on the input key. It is almost impossible to break the encryption without knowing the status of the key matrix which is randomized. The encrypted patterns shown above shows how different are the cipher text even when applied on similar text or the same text with subtle difference under same key. This shows the strength of the encryption. This method is applicable to different type of fields such as banking sectors, defence, short message encryption as well as large text encryption.

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Dr. Asoke Nath is Associate Professor in Department of Computer Science, St. Xavier's College (Autonomous), Kolkata-16. His major research areas are Cryptography and network security, Data hiding, Image processing, Green computing, E-learning.


Payel Pal, B.Sc (Hons). in Computer Science From University of Calcutta. Pesently pursuing MCA under IGNOU.


[^0]:    Asoke Nath, Department of Computer Science, St. Xavier's College(Autonomous), Kolkata, India.

    Payel Pal, Indira Gandhi National Open University,St. Xavier's College(Autonomous),Kolkata, India.

[^1]:    Bitwise Columnar transposition decryption function algorithm:col_trans_dec(file f1,file f2)

    This algorithm is the reverse process of col_trans_enc algorithm

