# Modern Encryption Standard version IV: (MES-IV)

# Asoke Nath<sup>1</sup>, Payel Pal<sup>2</sup>

#### Abstract

In the present paper the authors have introduced a new cryptographic method named as Modern Encryption System Standard version IV which is basically a symmetric key cryptographic method. Here the authors have used three different type of cryptographic methods .Those method are columnar transposition method, bit level generalized vernam cipher method with feedback and bit wise XOR operation. This system is the extension of MES-III and partly Bit level Encryption Standard (BLES) -II and III. BLES-I, II, III and MES-I, II, III developed by Nath et al where MES-I, II, III mostly based on mainly byte level encryption method. BLES-I, II and III are based on mostly bit level encryption methods. In the present MES-IV method authors have tried to combine different bit level encryption to make the entire encryption system more secured. The introduction of feedback in bit level generalized vernam cipher method prevents attacks such as differential attack or plain text attack. The random key generator has been used to construct the keypad for vernam cipher method. The present encryption method is free from common attacks and it is almost impossible to break the present method without knowing the exact key and the methods. This encryption method will be used to encrypt password, short messages, financial data etc.

#### Keywords

Plain text, Cipher text, Randomization, Columnar Transposition, Feedback

#### 1. Introduction

Keeping secrets is not easy. In fact human tendency is such that when told that something is a secret and asked to keep it secret, people are actually quite eager to share that secret with everyone else. In the early days of serious computing there was not a great deal of emphasis on security, because the systems in those days were proprietary or closed. The chances of someone getting an access to the information being exchanged were not very high.

As the minicomputers and microcomputers evolved in the 1970s and 1980s, the issue of information security started to gain more prominence. However it was the internet, which changed the whole computing paradigm and brought a tremendous change in the way computers communicated with each other. The world of computers had suddenly become very open. Therefore, it is very important to know how we can make information exchange secure. The confidentiality and genuineness of data has now become a very important issue. To send any important information from one user to another user normally the people are using e-mail as their transmission media. But the message of the e-mail can be trapped by the hacker between sender and receiver provided it is in raw form. To get rid of this problem one has to send the encrypted text or cipher text from client to server or to another client. To protect any kind of hacking problems nowadays network security and cryptography is an emerging research area where the programmers are trying to develop some strong encryption algorithm so that no intruder can intercept the encrypted key cryptography.

The cryptography methods can be divided into two categories: (i) symmetric key cryptography where one key is used for both encryption and decryption purpose. (ii) Public key cryptography where two different keys are used one for encryption and the other for decryption purpose. In symmetric key we have to maintain only one key and hence the key management is simple. In public key cryptography we maintain two keys one is public key which is known to everybody and that can be used for encryption purpose and there is another key called private key which is kept secret key and that is used for decryption purpose only.

Nath et al [1,2,3,4,5,6] developed different symmetric key cryptosystem. The advantage of symmetric key method is that key management is

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very simple. Recently Nath et al developed cryptography method called Modern Encryption Standard version-I and Modern Encryption Standard version-II. and Modern Encryption Standard version-III. The present version is called Modern Encryption Standard version-IV. Here the authors have used three independent cryptography methods namely columnar transposition method, bit-wise generalized vernam cipher method and bit wise XOR method. In the present version the different method can be used for large size of file encryption. The output shows that the encryption is very strong as the encrypted text is totally different if there is only one character different in two patterns. The standard encryption algorithm like RSA or DES if we apply on a pattern where all characters are same then after encryption the encrypted text will also show same repeated pattern. But the present method applied on repeated pattern but the output contains totally different pattern. In the present work the authors are proposing a symmetric key cryptography where firstly the keygen() function is called to generate the encryption number and the randomization number. Then the bit level columnar transposition encryption is applied block wise. After that bit level generalized vernam cipher method with feedback is applied. Hare the encryption key used is taken from the randomized two dimensional array. Finally the bit wise xor encryption is applied. This system of encryption can suitable to encrypt different financial secret data, corporate data, password, defense network etc.

#### 2. Algorithm

#### Bitwise Columner Transposition Encryption Function Algorithm : col trans enc(file f1,file f2)

# This function takes two files f1 and f2 as argument

step	1	:	Open f1 in read mode.
step	2	:	Open a file f2 in write mode.
step	3	:	set i=0 and n=0
step	4	:	set j=0
step	5	:	set mat[i][j]=n
step	6	:	set n=n+1
step	7	:	Set $j=j + 1$ and if $j < 16$ go to step 5
step	8	:	set i=i+1 and if i<16 go to step 4
step	9	:	Call randomization()
step	10	:	set i=0 and set j=0

step	11	:	Go to the start of f1
step	12	:	read next 256 character from f1 and
1			assign number of available character
			to n and assign the read characters to
			a[256]
sten	13	•	if $n \le 0$ go to step 56
step	14		set 1=0
sten	15	:	if $1 \ge n$ go to step 22
step	16	:	Call function char to hit(all hinary)
step	10	·	//this assign each bit of ch to elements
			of array binary[8]
ston	17		So $k=0$
step	10	÷	$f_{\rm L} = 0$
step	10	:	If $K \ge 0$ go to step 21 Set toble[1][1]_binemu[1] $\ge 49$
step	19	•	Set table[1][K]=Dinary[K] + 48
step	20	:	Set $k=k+1$ and go to step 18
step	21	:	Set I=I+1 and go to step 15
step	22	:	assign all element of filearr[8] to 0
step	23	:	set $11=0$ and $11=0$
step	24	:	set k=0
step	25	:	If $k \ge 8$ go to step 44
step	26	:	Set temp=mat[i][j]%8 + 1 // a%b
			returns the remainder after dividing a
			by b
step	27	:	set i=i+1
step	28	:	If $i=16$ set $i=0$ and $j=j+1$
step	29	:	If $j=16$ set $j=0$ .
step	30	:	set m=0
step	31	:	if m>=k go to step 34
step	32	:	if temp=filearr[m] go to step 34
step	33	:	increase m by 1 and go to step 31
step	34	:	if m <k 26<="" go="" step="" td="" to=""></k>
step	35	:	set filearr[k]=temp
step	36	:	Set $k=k+1$
step	37	:	set 1=0
step	38	:	if l>=n go to step 43
step	39	:	set table1[i1][j1]=table[1][temp-1]
step	40	:	set $j_{1=j_{1+1}}$
step	41	:	if $i1=8$ set $i1=0$ and $i1=i1+1$
step	42	:	set l=l+1 and go to step 38
step	43	:	Go to step 25
step	44	:	set 1=0
step	45	:	if l>=n go t step 53
step	46	:	set add=0
step	47	:	set k=0
step	48	•	if $k \ge 8$ go to step 51
step	49	÷	set add=add+ (table1[1][k]-
•P	. /	•	48)*power(7-i)
sten	50	•	set $k=k+1$ and go to step 48
sten	51	:	Write add to file f2
sten	52	÷	set $l=l+1$ and go to step 45
sten	53	:	if $n < 256$ go to step 56
sten	54	:	read next 256 character from f1 and
step	51	•	assign number of available character

			to n and assign the read characters to	step	42	:	increase k by 1
			a[256]	step	43	:	increase j by 1 and go to step 40
step	55	:	Go to step 13	step	44	:	set $i=i+1$ and go to step 38
step	56	:	Close all the files.	step	45	:	call randomization()
step	57	:	stop	step	46	:	set i=j=0
			-	step	47	:	set cr1=0
			bitwise vernam encryption with	step	48	:	read next character from f1 and
			feedback:vernambitenc(file				assign to ch
			input,file output)	step	49	:	if eof is found go to step 70
step	1	:	set k=0	step	50	:	call char_to_bit(ch,bitpattern[8])
step	2	:	set i=0	step	51	:	call char_to_bit(mat[i][j],key_bit[8])
step	3	:	if $i \ge 16$ go to step 10	step	52	:	set i=i+1
step	4	:	set j=0	step	53	:	if $i=16$ set $i=0$ and set $j=j+1$
step	5	:	if $j \ge 16$ go to step 9	step	54	:	if j=16 set j=0
step	6	:	set mat[i][i]=k	step	55	:	set
step	7	:	increase k by 1	1			cr=(bitpattern[0]+key bit[0]+cr1)%2
step	8	:	increase j by 1 and go to step 5	step	56	:	set cb[0]=cr1=cr
step	9	:	set $i=i+1$ and go to step 3	step	57	:	set k=1
step	10	:	call randomization()	step	58	:	if $k \ge 8$ go to step 63
step	11	•	set i=i=0	step	59	:	set
step	12	•	set cr1=0				cr = (bitpattern[k] + kev bit[k] + cr1)%2
step	13	•	open input file as fpn	step	60	:	set cb[k]=cr
step	14	•	read next character from fpn and	step	61	•	set cr1=cr
~~··r			assign to ch	step	62		set $k=k+1$ and go to step 58
step	15	·	if eof is found go to step 36	step	63	÷	set add=0
step	16	•	call char to bit(ch.bitpattern[8])	step	64	•	set k=0
step	17	•	call char to bit(mat[i][i].kev bit[8])	step	65		if $k \ge 8$ go to step 68
step	18	•	set $i=i+1$	step	66	•	set add=add+cb[k]*power(7-k)
step	19		if $i=16$ set $i=0$ and set $i=i+1$	step	67		increase k by 1 and go to step 65
step	20	:	if $i=16$ set $i=0$	step	68		write add to file f?
step	21	:	set	step	69		go to step 48
step	-1	•	cr=(bitnattern[0]+key bit[0]+cr1)%2	sten	70	:	set $k=0$
sten	22		set $ch[0]=cr1=cr$	sten	71	:	set i=0
sten	22	:	set $k-1$	sten	72	:	if $i > -16$ go to step 79
sten	23	:	if $k > -8$ go to step 29	sten	73	:	set $i=0$
sten	25	:	set	sten	74	:	if $i > -16$ go to step 78
step	25	•	cr = (hitnattern[k] + key hit[k] + cr1)%2	sten	75	:	set mat[i][i]=k
sten	26		set $ch[k] - cr$	sten	76	:	increase k by 1
step	20	:	set cr1-cr	sten	70	:	increase i by 1 and go to step 5
step	27	:	set $k = k + 1$ and go to step 24	step	78	:	set $i=i+1$ and go to step 72
step	20	:	set $\mathbf{x} = \mathbf{x} + 1$ and go to step 24	step	70	:	call randomization()
stop	29	•	set $k=0$	step	80	:	sot i-i-0
step	31	:	if $k = 8$ go to step 34	step	80 81	:	set $r_{-}=0$
step	22	:	$\lim_{k \to -\infty} \sum_{i=1}^{n} \sum_{j=1}^{n} \sum_{i=1}^{n} \sum_{i=1}^{n} \sum_{i=1}^{n} \sum_{i=$	step	01 07	:	read next character from f2 and
step	32 22	:	set add=add+co[K] power(/-K)	step	02	•	read next character from 12 and
step	23 24	•	write add to file fl	aton	02		assign to ch
step	24 25	:	while add to me m	step	03	:	all abor to hit(ab hitpattam[9])
step	20	÷		step	04 05	•	call char_to_bit(cn,bitpattern[8])
step	20 27	÷	Set $K = 0$	step	0J 02	:	$can char_{10}_{011}(hat[1][1], key_01t[8])$
step	21 20	÷	Still $I = 0$ if $i > -16$ go to star 45	step	00 07	:	Set $1-1+1$ if $i-16$ set $i-0$ and set $i-1$
step	20 20	÷	$\frac{11}{12} = 10 \text{ go to step 45}$	step	0/ 00	:	If $i=16$ set $i=0$ and set $j=j+1$
step	39 40	:	Set $J=0$ if $i > -16$ go to step 44	step	00 00	:	II $j=10$ set $j=0$
step	4U 41	:	II $J \ge 10$ go to step 44	step	69	:	
step	41	:	set $mat[1][]=K$				$cr=(oupattern[0]+key_bit[0]+cr1)%2$

step	90	:	set cb[0]=cr1=cr	step	139
step	91	:	set k=1	step	140
step	92	:	if $k \ge 8$ go to step 97	_	
step	93	:	set		
			$cr = (bitpattern[k] + key_bit[k] + cr1)\%2$		
step	94	:	set cb[k]=cr		
step	95	:	set cr1=cr		
step	96	:	set $k=k+1$ and go to step 92		
step	97	:	set add=0		
step	98	:	set k=0		
step	99	:	if k>=8 go to step 102		
step	100	:	set add=add+cb[k]*power(7-k)	stan	1
step	101	:	increase k by 1 and go to step 99	step	2
step	102	:	write add to file f3	step	3
step	103	:	go to step 82	step	1
step	104	:	call file rev(f3,f4)	step	-
step	105	:	set k=0	step	5
step	106	:	set i=0	ston	6
step	107	:	if $i \ge 16$ go to step 114	step	0
step	108	÷	set i=0	step	0
step	109	:	if $i \ge 16$ go to step 113	step	0
step	110	÷	set mat[i][i]=k	step	9
step	111	•	increase k by 1	step	10
step	112	:	increase i by 1 and go to step 109	step	11
step	112	:	set $i=i+1$ and go to step 107	step	12
step	114	:	call randomization()	step	13
sten	115	:	set $i=i=0$	step	14
sten	116	:	set $cr1-0$	step	15
sten	117	:	read next character from f/ and	step	16
step	11/	·	assign to ch	step	17
sten	118		if eof is found go to step 139	step	18
step	110	:	call char to hit(ch hitpattern[8])	step	19
stop	120	:	call char to bit(mat[i][i] kay bit[8])		
step	120	:	$can char_{0} bh(hat[1][1], key_bh[0])$	step	20
step	121	:	Set $I = I + I$ if $i = 16$ set $i = 0$ and set $i = i + 1$	step	21
step	122	•	If $I=10$ set $I=0$ and set $J=J+1$	step	22
step	123	÷	$\lim_{n \to \infty} j = 10 \text{ set } j = 0$	step	23
step	124	:		step	24
	105		$cr=(bitpattern[0]+key_bit[0]+cr1)%2$	step	25
step	125	:	set cb[0]=cr1=cr	step	26
step	126	:	set k=1		
step	127	:	If $k \ge 8$ go to step 132	step	27
step	128	:	set		
	1.00		$cr=(bitpattern[k]+key_bit[k]+cr1)%2$	step	28
step	129	:	set cb[k]=cr	step	29
step	130	:	set cr1=cr	step	30
step	131	:	set k=k+1 and go to step 127	-	
step	132	:	set add=0		
step	133	:	set k=0	Bity	wise (
step	134	:	if $k \ge 8$ go to step 137	algo	orith
step	135	:	set add=add+cb[k]*power(7-k)		
step	136	:	increase k by 1 and go to step 134	Thi	s alor
step	137	:	write add to file output file	aloc	rithn
step	138	:	go to step 117	aige	/111111

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**Bit Wise XOR Encryption** 

: close all files

: stop

# Function: bitxorenc(File F1,File **F2**)

#### This function takes two files f1 and f2 as argument

step	1	:	Open the file f1 in read mode	
step	2	:	Open the file f2 in write mode	
step	3	:	set l=size of file f1	
step	4	:	set n1=1/32	
step	5	:	set n1=l%32 // a%b returns the	
			remainder after dividing a by b	
step	6	:	Go to the start of file f1	
step	7	:	set i=0 and n=0	
step	8	:	set j=0	
step	9	:	set mat[i][j]=n	
step	10	:	set n=n+1	
step	11	:	Set $j=j + 1$ and if $j<16$ go to step 9	
step	12	:	set i=i+1 and if i<16 go to step 8	
step	13	:	set i=1	
step	14	:	if i>secure then go to step 17	
step	15	:	call randomization()	
step	16	:	set i=i+1 and go to step 14	
step	17	:	set i=1	
step	18	:	if i>n1 then go to step 23	
step	19	:	Read next 32 character from file f1	
			and assign it to array data1[32]	
step	20	:	Call bit_stream(data1[32])	
step	21	:	Call encrypt_bit()	
step	22	:	set i=i+1 go to step 18	
step	23	:	if n2=0 go to step 30	
step	24	:	set i=0	
step	25	:	if $i \ge n2$ go to step 30	
step	26	:	Read next character fron file f1 and	
			assign to data2[i] of array data2[32]	
step	27	:	set	
			data2[i]=rshift_residual(data2[i],5)	
step	28	:	write data2[i] to file f2	
step	29	:	set $i=i+1$ go to step 25	
sten	30	•	close all files	

#### Columnar transposition decryption function m:col\_trans\_dec(file f1,file f2)

orithm is the reverse process of col\_trans\_enc n

#### bitwise vernam decryption with feedback:vernambitdec(file input,file output)

This algorithm is the reverse of vernambitenc algorithm.

#### Bitwise Xor Decryption Function: Bitxordec(File F1,File F2)

This algorithm is the reverse of bitxorenc algorithm.

## 3. Results and Discussion

The Modern Encryption Standard version IV (MES-IV) is applied on different type of text. This is applied on various repeated patterns such as 1024 numbers of ASCII 0 to ASCII 16. The corresponding frequency graphs are shown below. The graph shows the diverse nature of encrypted text even the input text is of repeated same character.





#### Fig-1: Frequency Graph of ASCII code '0','1','2','3'

This encryption method applied to different type of text/patterns and shown below are the pairs of such different type of patterns and the corresponding cipher text.

#### Table-1: Some known plain texts vs. encrypted text

S1.	ORIGINAL TEXT	ENCRYPTED TEXT
No		
1	АААААААААААААА	薪 <b>갓血 必</b> 溷皤 <b>鼀黯晰浤</b>
	АААААААААААААА	
	АААААААААААААА	<b>됟嚿.</b> 稠砲□酘에 <b>脆黺</b> 稲
	АААААААААААААА	<b>펖</b> 2ど音娟銀版
	AAAAAAAA (	
	64-As)	
2	BAAAAAAAAAAAAAA	ÇŸiKw à-T∙ô~ hed]»k Ôï
	АААААААААААААА	i, žfö r
	АААААААААААААА	í©VôYÈÿ6®Ô°¦ÖïM£ÛuÎâ
	АААААААААААААА	á• AÔì J£• u ðà™/
	АААААААА	
	(B + 64 As)	

3	AAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA	Ç>->o mTŸð• èAd]«b ÔïD⁻¶ôr i©VôYÈÿ6®Ô°¦ÖïM£ÛuÎâ á•AÔì J£•u ðà™⁄
4	0000000	J i!′ý
5	1111111	" J i!′ý
6	01010101	– J i!´ý
7	1111111100000000	[0A¢RZ£ WÜ¢R ‡p
8	HE IS GOOD	¥ÒWÎê°Œ£5w
9	abcabcabc	"• Då¢x° 5
10	HE IS GOON	¥Ðuìê°Œ£5w
11	CE IS GOON	′Àq¨®°Œ£5w

In the table below the plain text and the corresponding encrypted text are shown. The text of Sl. No. 12 & 13 are exactly same except the fourth character. However the encrypted texts are quite different under the same encryption key. The Sl. No. 14 & 15 shows the same thing.

	ORIGINAL	ENCRYPTED TEXT
	TEXT	
12	The root of all	Ç8
	modern	iØãÕ∙ y=g' -
	programming	C e=f@Eal!2BNWE¢O§•1/4e
	languages is	¨i>Q>³ùÚ Çcβ
	ALGOL, introduc	Û• 'VNTÁ6Cçæ» Á
	ed in the early	[Â.RŽ3Ï′¬Œ'µþ ðfK-
	1960s.ALGOL	pE¦EŽbÒŸ'ž•GÉ¥Fã
	was the first	s7'aN€@ ìÉq kú• °• . H~Đ4
	computer	• ác¼»¶ <sup>™</sup> !"Ws®ðXäÅž
	language to use a	Ú;:V^œ•CY8ñ³t6%sγ•
	block structure.	%èF¤4cz <sup>2</sup> elmfE…ó
	Although it never	êf‡uÙu '»::
	became popular	$\tilde{N}aaa \otimes \hat{I} + \hat{a} \bullet \hat{a} \boxtimes \hat{O} / \ddot{A} \check{S} \times \langle \hat{O} \tilde{O} \rangle$
	in USA, it was	10W fv -
	widely used in	#●(ŠÉÙ·Ü(©×«Èšm>s
	Europe. ALGOL	$\neg RUII fnR$
	gave the concept	őcfiÑDŠ«ZKE@Äà}z ∐>Ú
	of structured	«h
	programming to	•½ÈÀ€!Cìù4 Æêð‰» a"[e
	the computer	
	science	
	community.	

13	The boot of all	iæ(›kÄ ÔÁy<æÓ
	modern	ø• ý%å• AÕ£
	programming	;7BÅwR_tC®¬~ ′ïh~ P
	languages is	$\partial \ddot{\mathbf{F}} \bullet \mathbf{M} \mathbf{F} # \mathbf{I} \mathbf{I} \mathbf{g} & \mathbf{I} \mathbf{O} \mathbf{A}^{"c} \mathbf{A} = \mathbf{r} \mathbf{a}$
	ALGOL introduc	$\int \frac{1}{\sqrt{23}} dx = \frac{1}{\sqrt{23}} \int $
	ed in the early	$1A2^{4}$ SIµ-1
		$a_{1}$ $\frac{1}{\sqrt{41}}$ $\frac{1}{\sqrt{405}}$ $\frac{1}{\sqrt{405}}$ $\frac{1}{\sqrt{205}}$
	1900S.ALOOL	$\pm <$ 1fU3Dyerw TArV"H1‡1u
	was the first	SE
	computer	. häð−{‡• wå W·Ø!•Ö2œ¥
	language to use a	e…F Û/{& ™•Óy¢Ñ 5åÌ⁻a
	block structure.	ÉèG%uO/V3%ogÕÅ $f$ >ï $f$ "
	Although it never	•C•iÑ›·*
	became popular	$Aaaassi \pm \delta $
	in USA, it was	IOMEve
	widely used in	OMLY -
	Europe, ALGOL	#●(SEU·U(©×«Esm>s
	gave the concept	¬RU![fpR
	of structured	õcfiNDS«ZKE@Aà}z U>U
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	into a timesharing	<sup>1</sup> ⁄₂• H−gú Zå±—£;
	system. In this	'J)X uHg a¨; Ob⁻KÕ>ÍÏ
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	the data are typed	$\hat{\partial} \hat{\partial} \hat{\partial} \hat{\partial} \hat{d} \hat{d} \hat{d} \hat{d} \hat{d} \hat{d} \hat{d} d$
	into the computer	$(0)^{\circ} (0)^{\circ} (0)^$
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	terminal or a	I3ìUZOœ∙äUK%°
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	computer.	
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	generally tied	ÖC÷.6_¿~É×U⁻Hû¥j
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	the data are typed	$a_{a} M \delta E^{2} N - ch v_{a} 1 \mu k$
	into the computer	• $\hat{\lambda}$ $\cdot 0 \in \mathbb{N}$ $D = \hat{\lambda} = 0$
	via a timesharing	
	terminal or a	151UZU@ aUK%~
	norsonal	
	personal	
1	computer.	1

# 4. Conclusion

The results above show how strong the encryption method is. Three different methods of encryption

techniques have been applied. The feedback mechanism gives to the strength of the encryption. The random matrix is generated on the input key. It is almost impossible to break the encryption without knowing the status of the key matrix which is randomized. The encrypted patterns shown above shows how different are the cipher text even when applied on similar text or the same text with subtle difference under same key. This shows the strength of the encryption. This method is applicable to different type of fields such as banking sectors, defence, short message encryption as well as large text encryption.

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